

$[n+1] = \{0, 1, \dots, n\}$   
 $[n] = \{0, 1, \dots, n-1\}$

Recall:  $\{1, 2\}$   $\{1, 2, 3\}$  The formula with Gödel number  $\text{sub}(y, y, 1)$  is not provable.

$A(n, p) = \begin{cases} 1 & \text{if } \phi(p) = n \\ 0 & \text{otherwise} \end{cases}$

$\phi(A) = n \iff$  Gödel number.

$n$ : numeral, encoding of a formula.

$\text{proof}(n, m) = \begin{cases} 1, & \text{if } n = \phi(\text{proof of } p) \text{ s.t. } \phi(p) = m, \text{ "n is a proof of m"} \\ 0 & \text{otherwise} \end{cases}$

$\text{THM}(b) := (\exists x) (\text{proof}(x, b))$

"b has a proof"

$\exists x \text{ s.t. } x = \phi(\text{proof of } p) \text{ and } \phi(p) = b.$

$\text{sub}(m, 1) = \phi(\exists x (x = Sx))$

$a_{n+2} = \frac{3a_{n+1}}{3(3a_n - 2a_n)}$

$\begin{pmatrix} 7 & -6 \\ 3 & -2 \end{pmatrix}$

$\begin{pmatrix} a_{n+2} \\ a_{n+1} \end{pmatrix}$

Want refer this to itself.

$\neg \text{THM}(b)$  "b is not provable"

Define a function.  $\text{sub}(m, n) = j$  if.

$m = \phi(A(x))$  where  $x$  is the only free variable.

$j = \phi(A(n))$  ( $n$ ) numeral of  $n$ .

$\text{sub}(m, n)$  does:

- compute  $m$  to get  $A(x)$
- take  $n$  as a number.
- do  $A(n)$
- compute  $j = \phi(A(n))$

upshot:  $\text{sub}(m, n)$  is computable over  $\mathbb{N}$ , it is an arithmetic function.

Gödel has to define  $\text{SUB}(x, y, z)$  as an explicit formula.

Consider  $A(x) := \neg \text{THM}(\text{sub}(x, x))$

$\phi(A(x)) = \underline{n}$

$G = A(\underline{n})$  is not provable & true.

If  $G$  is not true by  $A(x)$  is possible.

If  $G$  is provable by  $A(x)$  it is not true.

it self has a formula in the formal system.