

primitive propositions  
↓  
compound statements

$\neg$   $\vee$   $\wedge$   $\rightarrow$   $\forall$   $\exists$   
 $A^c$   $A \cup B$   $A \cap B$   $A \setminus B$   
 logic set.

$U$  ambient set

$\{x \in U \mid \neg P(x)\}$   
 property.  $U = \mathbb{R}$   
 $\{y \in \mathbb{R} \mid y > 2\}$   
 extension of  $\neg P(x)$  in  $U$ .

properties	
$P(x)$	$\{x \in U \mid P(x)\}$
$P(x)$ and $Q(x)$	$\{x \in U \mid P(x)\} \cap \{x \in U \mid Q(x)\}$
or	
$\neg P(x)$	$\{x \in U \mid P(x)\}^c \rightarrow \text{comp in } U$
$Q(x) \Rightarrow P(x)$	$\{x \in U \mid Q(x)\} \subseteq \{x \in U \mid P(x)\}$
$\neg(Q(x) \vee P(x))$	$\{x \in U \mid Q(x)\}^c \cap \{x \in U \mid P(x)\}^c$

✓ First order logic  $\phi(x) \in \{ \neg, \wedge, \vee \}$   
 quantifier on  $x$ .

Second order logic = FOL

A foundation of math

+ quantifiers on predicates.

Model  $\rightarrow M$

universe of states  $U_M$

For any proposition  $A$ , we associate an interpretation  $\llbracket A \rrbracket_M \subseteq U_M$ ,

s.t.  $\llbracket \top \rrbracket_M = U_M$   
 $\llbracket \perp \rrbracket_M = \emptyset$

$\llbracket A \wedge B \rrbracket_M = \llbracket A \rrbracket_M \cap \llbracket B \rrbracket_M$   
 $\llbracket A \vee B \rrbracket_M = \llbracket A \rrbracket_M \cup \llbracket B \rrbracket_M$



$$\begin{aligned}
 \llbracket A \vee B \rrbracket_M &= \llbracket A \rrbracket_M \cup \llbracket B \rrbracket_M \\
 \llbracket \neg A \rrbracket_M &= \llbracket A \rrbracket_M^c
 \end{aligned}$$

S model  $U_S \rightarrow$  students

$$\llbracket a \rrbracket_S = \{x \in U_S \mid x \text{ is master major}\}$$

$$\llbracket b \rrbracket_S = \{x \in U_S \mid x \text{ has a driver licence}\}$$

$$\llbracket c \rrbracket_S = \{x \in U_S \mid x \text{ is nineteen}\}$$

Def: Proposition  $A$ ,  $A$  is **valid** in  $\mathcal{M}$

$$\downarrow \text{ iff. } \llbracket A \rrbracket_{\mathcal{M}} = U_{\mathcal{M}}$$

prop  $A, B$   $A$  **entails**  $B$  in  $\mathcal{M}$

$$\text{iff. } \llbracket A \rrbracket_{\mathcal{M}} \subseteq \llbracket B \rrbracket_{\mathcal{M}}$$

prop  $A$  is valid iff  $\neg A$  is valid in all models

$$\vDash \neg \neg A$$

prop  $A, B$ ,  $A$  entails  $B$  iff.  $A$  entails  $B$  in all models

$$\vdash A \rightarrow B$$

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1/s

1/s



e/s

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11 specific model - truth assignment.

